

Charles University in Prague Faculty of Mathematics and Physics

Suffix tree construction with minimized branching MASTER'S THESIS

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September 3, 2012





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Suffix tree

A tree-like data structure for performing fast search-like operations on strings.

Properties

- Search for a pattern occurrence in $\mathcal{O}(|pattern|)$ time.
- Most implementations require 20–36 bytes per each input character in the worst case.

Some applications

- Bioinformatics
 - searching for patterns in DNA and protein sequences
- Finding repetitive text structures
- Pattern matching using wildcards or regular expressions



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Example suffix tree on top of the text ABABBA\$







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Introduction

Suffix tree construction

- entirely in memory
- over a sliding window





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Goals

- theoretical analysis
- implementation
- experimental evaluation



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Minimized branching

What is it?

An alternative method for suffix link simulation introduced by Senft and Dvořák, 2012.





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Suffix link simulation

- Top-down
- Bottom-up



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Applicable to the algorithms by

- McCreight (1976)
- Ukkonen (1992)



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Examined methods of suffix tree construction

Algorithms

- McCreight's (1976)
- Ukkonen's (1992)
- Partition and Write Only Top Down PWOTD (Tata, Hankins, Patel, 2004)



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Implementation techniques

- Simple Linked List (Kurtz, 1999)
- Simple Hash Table (Kurtz, 1999)
- Simple Linear Array (Tata, Hankins, Patel, 2004)



Experiments conducted

Implementations used

- implementation by Senft and Dvořák, 2012
- SLLI implementation by S. Kurtz, 1999
- PWOTD implementation by Tata, Hankins, Patel, 2004
- our own implementation



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Data used

- pseudorandom input files
- standard corpus files (Pizza & Chili and Lightweight corpus)
- special input files



Description of the algorithms

• We have presented unified descriptions and definitions of every algorithm and its variation analyzed in this thesis.





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The implementation

- Every algorithm and implementation technique is implemented using similar level of detail and quality.
- Compilable on typical UNIX platforms.
 - $\bullet\,$ tested on Linux and Mac OS X



Usage recommendations

- McCreight's / Ukkonen's algorithms vs. PWOTD
 - PWOTD often computes the length of the lcp
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 - higher alphabet size \implies use hash table



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- linked lists vs. hash table
 - lower alphabet size \implies use linked lists
 - $\bullet\,$ higher alphabet size $\implies\,$ use hash table
- percolating update vs. batch update
 - both methods have constant amortized time complexity
 - percolating update makes use of suffix tree traversal during the construction \implies its constant is smaller



Experimental results

McCreight's / Ukkonen's algorithms vs. PWOTD

- PWOTD is faster on pseudorandom files
- McCreight's / Ukkonen's algorithms are faster on corpus files
 - PWOTD is slow on files with large average lcp





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top-down vs. bottom-up suffix link simulation

- bottom-up method usually outperforms top-down
- suffix tree construction time is reduced by 5-10%
 - except for special, adversary strings



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- suffix tree construction time is reduced by 5–10%
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percolating update vs. batch update

- percolating update is *almost* always faster
 - $\bullet\,$ pseudorandom files: construction time is reduced by 5–10%
 - $\bullet\,$ corpus files: the reduction is larger, typically around $15\%\,$



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Suffix tree traversal log

Human readable suffix tree representation: P(number)[depth]--"label"(length)-->C(number)[depth]{suffix link}



DFS traversal log

```
Suffix tree traversal BEGIN
P(1)[0] -- "a"(1) --> C(5)[1]{1}
P(5)[1] - "b"(1) - >C(2)[2]{3}
P(2)[2] -- "abba$"(5) --> C(-1)[7]
P(2)[2] - - ba$"(3) - > C(-3)[5]
P(5)[1]--"$"(1)-->C(-6)[2]
P(1)[0]--"b"(1)-->C(3)[1]{1}
P(3)[1] - - a''(1) - > C(4)[2]{5}
P(4)[2] -- "bba$"(4) -->C(-2)[6]
P(4)[2] - "$"(1) - > C(-5)[3]
P(3)[1]--"ba$"(3)-->C(-4)[4]
P(1)[0]--"$"(1)-->C(-7)[1]
Suffix tree traversal END
```



Repetitive structures

maximal pair

a pair of identical substrings impossible to extend in any direction $\dots GTTATTATT_{L}^{R}TTATT_{R}^{R}$...

maximal repeat

a substring whose occurrences form a maximal pair ... T<u>AC</u>TGAC<u>GT</u>TGTC...

supermaximal repeat

A maximal repeat which is not a substring of any other maximal repeat. In the text ACGCCCCTACGA:

- CG is a maximal repeat, but it is not supermaximal
- ACG is a supermaximal repeat